Matrix Multiplicative Weights and Nearly-Linear Time Robust Statistics

Kevin Tian (UT Austin) New Frontiers in Robust Statistics Workshop TTIC, 2024

with many thanks to...





...and more

Talk outline

- A gentle introduction to MMW
 - Regret minimization
 - Matrix analysis
 - Implementation
 - Relatives of MMW
- Robust statistics primitives via MMW
 - Mean estimation
 - A tour of applications

References:

- MMW intro
 - Course notes (Continuous Algorithms, Spring '24)
 - Lectures 3, 5-7
 - Email me for pointers
- Robust statistics applications
 - ... Continuous Algorithms, Spring '25?
 - Email me for pointers

Regret minimization



Regret minimization

Goal: predict the future!

$$\sum_{t \in [T]} \ell_t(x_t) - \min_{x^* \in \mathcal{X}} \sum_{t \in [T]} \ell_t(x^*)$$

Importantly, x_t played before ℓ_t known (Also, ℓ_t can be adaptive)



Regret minimization

Goal: sublinear regret $\mathcal{X} \subset \mathbb{R}^d$ $\sum_{x^{\star} \in \mathcal{X}} \ell_t(x_t) - \min_{x^{\star} \in \mathcal{X}} \sum_{x^{\star} \in \mathcal{X}} \ell_t(x^{\star}) = o(T)$ $t \in [T]$ convex, compact $t \in [T]$ action space (e.g. norm ball) Importantly, x_t played before ℓ_t known (Also, ℓ_t can be adaptive)

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Linear regret minimization
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Goal: sublinear regret

$$\sum_{t \in [T]} \langle g_t, x_t \rangle - \min_{x^* \in \mathcal{X}} \sum_{t \in [T]} \langle g_t, x^* \rangle = o(T)$$
convex, action (e.g. not limportantly, x_t played before g_t known (Also, g_t can be adaptive)



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Linear regret minimization
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Application I: convex optimization

$$\sum_{t \in [T]} \langle g_t, x_t \rangle - \min_{x^* \in \mathcal{X}} \sum_{t \in [T]} \langle g_t, x^* \rangle = o(T) \qquad g_t = \nabla f(x_t) \qquad \text{adaptive!}$$

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Linear regret minimization
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Application I: convex optimization

$$\sum_{\substack{t \in [T] \\ Regret_T}} \langle g_t, x_t \rangle - \min_{x^\star \in \mathcal{X}} \sum_{t \in [T]} \langle g_t, x^\star \rangle = o(T) \qquad g_t = \nabla f(x_t) \quad \text{adaptive!}$$

$$Regret_T$$

$$Regret_T \ge \sum_{t \in [T]} f(x_t) - f(x^\star) \ge T(f(\bar{x}) - f(x^\star))$$

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Linear regret minimization
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Application I: convex optimization

$$\begin{split} \sum_{t \in [T]} \langle g_t, x_t \rangle &- \min_{x^\star \in \mathcal{X}} \sum_{t \in [T]} \langle g_t, x^\star \rangle = o(T) \qquad g_t = \nabla f(x_t) \quad \text{adaptive!} \\ \\ \hline \\ \text{Regret}_T \\ \\ \text{Regret}_T \geq \sum_{t \in [T]} f(x_t) - f(x^\star) \geq T(f(\bar{x}) - f(x^\star)) \\ & \text{vanishing suboptimality gap!} \\ & \text{+ works for cvx-ccv saddle point} \end{split}$$

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Linear regret minimization
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Application 2: dual certificates

$$\sum_{t \in [T]} \langle g_t, x_t \rangle - \min_{x^* \in \mathcal{X}} \sum_{t \in [T]} \langle g_t, x^* \rangle = o(T)$$

Regret_T

$$\mathcal{X} := \{ x \mid \|x\| \le 1 \}$$

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Linear regret minimization
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Application 2: dual certificates

$$\sum_{t \in [T]} \langle g_t, x_t \rangle - \min_{x^* \in \mathcal{X}} \sum_{t \in [T]} \langle g_t, x^* \rangle = o(T) \qquad \qquad \mathcal{X} := \{ x \mid ||x|| \le 1 \}$$
Regret_T

$$-\min_{x\in\mathcal{X}}\sum_{t\in[T]}\langle -g_t, x^*\rangle = \max_{x\in\mathcal{X}}\left\langle\sum_{t\in[T]}g_t, x^*\right\rangle = \left\|\sum_{t\in[T]}g_t\right\|_*$$

...if we choose g_t , regret minimization algos certify bounds

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Linear regret minimization
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Linear regret minimization
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(probability) simplex

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Linear regret minimization
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Examples:

Application 2: dual certificates

$$\begin{array}{ll} \operatorname{Action \, set} & \max_{\mathbf{X}^{\star} \in \mathcal{X}} \langle \mathbf{S}, \mathbf{X}^{\star} \rangle \\ \mathcal{X} = \left\{ \mathbf{X} \in \operatorname{Sym}^{d \times d} | \, \|\mathbf{X}\|_{q} \leq 1 \right\} & \|\mathbf{S}\|_{p}, \, \frac{1}{p} + \frac{1}{q} = 1 \\ & \text{Schatten } q\text{-norm ball} \\ \mathcal{X} = \left\{ \mathbf{X} \in \operatorname{PSD}^{d \times d} | \, \operatorname{Tr}(\mathbf{X}) = 1 \right\} & \lambda_{\max}(\mathbf{S}) \end{array}$$

Spectraplex

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Linear regret minimization
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Examples: **Application 2: dual certificates** $\max_{\mathbf{X}^{\star}\in\mathcal{X}}\langle\mathbf{S},\mathbf{X}^{\star}\rangle$ Action set $\mathcal{X} := \{ x \mid ||x|| \le 1 \}$ $\mathcal{X} = \left\{ \mathbf{X} \in \operatorname{Sym}^{d \times d} \mid \left\| \mathbf{X} \right\|_{q} \le 1 \right\} \quad \left\| \mathbf{S} \right\|_{p}, \ \frac{1}{n} + \frac{1}{q} = 1$ von Neumann trace inequality: Schatten q-norm ball $\max \langle \mathbf{V} \mathbf{D}_2 \mathbf{V}^\top, \mathbf{U} \mathbf{D}_1 \mathbf{U}^\top \rangle$ V unitary $\mathcal{X} = \{ \mathbf{X} \in \text{PSD}^{d \times d} \mid \text{Tr}(\mathbf{X}) = 1 \} \quad \lambda_{\max}(\mathbf{S})$ achieved iff $\mathbf{V} = \mathbf{U}$ up to permutation and subspace invariance Spectraplex

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Linear regret minimization
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Application 2: dual certificates

$$\begin{array}{ll} \operatorname{Action \, set} & \max_{\mathbf{X}^{\star} \in \mathcal{X}} \langle \mathbf{S}, \mathbf{X}^{\star} \rangle \\ \\ \mathcal{X} = \left\{ \mathbf{X} \in \operatorname{Sym}^{d \times d} \mid \|\mathbf{X}\|_q \leq 1 \right\} & \|\mathbf{S}\|_p, \ \frac{1}{p} + \frac{1}{q} = 1 \\ \\ & \text{Schatten } q\text{-norm ball} \\ \\ \\ \mathcal{X} = \left\{ \mathbf{X} \in \operatorname{PSD}^{d \times d} \mid \operatorname{Tr}(\mathbf{X}) = 1 \right\} & \lambda_{\max}(\mathbf{S}) \end{array}$$

$$\mathcal{X} := \{ x \mid \|x\| \le 1 \}$$

Upside:

Most SOTA fast robust stats algos based on this connection to regret minimization!

Spectraplex

Linear regret minimization

Examples:

Application 3: SDP feasibility

$$\mathcal{X} = \left\{ \mathbf{X} \in \mathrm{PSD}^{d \times d} \mid \mathrm{Tr}(\mathbf{X}) = 1 \right\}$$

Spectraplex

$$\exists y \in \mathcal{Y} : \sum_{i \in [n]} y_i \mathbf{A}_i \succeq \mathbf{0}_d?$$

Linear regret minimization

Examples:

Application 3: SDP feasibility

$$\mathcal{X} = \left\{ \mathbf{X} \in \mathrm{PSD}^{d \times d} \mid \mathrm{Tr}(\mathbf{X}) = 1 \right\}$$

Spectraplex

$$\exists y \in \mathcal{Y} : \sum_{i \in [n]} y_i \mathbf{A}_i \succeq \mathbf{0}_d?$$

Regret minimization for approximate saddle point of:

$$\min_{\mathbf{X}\in\mathcal{X}}\max_{y\in\mathcal{Y}}\left\langle \mathbf{X},\sum_{i\in[n]}y_i\mathbf{A}_i\right\rangle$$





 $v^{\top} \nabla^2 \varphi(x) v \ge \|v\|^2$

$$\|g_t\|_* \leq G \text{ for all } t$$
$$\varphi : \mathcal{X} \to \mathbb{R} \text{ is 1-s.c. in } \|\cdot\|$$
$$\max_{x \in \mathcal{X}} \varphi(x) - \min_{x \in \mathcal{X}} \varphi(x) \leq \Theta$$

$$\max_{x^{\star} \in \mathcal{X}} \sum_{t \in [T]} \langle g_t, x_t - x^{\star} \rangle \lesssim G\sqrt{\Theta T}$$

Different regularity: RHS is $poly(G, \Theta)$

$$\|g_t\|_* \leq G \text{ for all } t$$
$$\varphi : \mathcal{X} \to \mathbb{R} \text{ is 1-s.c. in } \|\cdot\|$$
$$\max_{x \in \mathcal{X}} \varphi(x) - \min_{x \in \mathcal{X}} \varphi(x) \leq \Theta$$

$$\max_{x^{\star} \in \mathcal{X}} \sum_{t \in [T]} \langle g_t, x_t - x^{\star} \rangle \lesssim G \sqrt{\Theta T}$$

Aside: convex duality

$$\varphi^*(y) := \max_{x \in \mathcal{X}} \langle y, x \rangle - \varphi(x)$$

Conjugate of convex function



$$\varphi^*(y) := \max_{x \in \mathcal{X}} \langle y, x \rangle - \varphi(x) \qquad \begin{array}{c} \text{Conjugate of} \\ \text{convex function} \end{array}$$

$$\nabla \varphi^*(y) = \mathrm{argmax}_{x \in \mathcal{X}} \langle y, x \rangle - \varphi(x) \qquad \begin{array}{l} \mathrm{Maximizing} \\ \mathrm{argument} \end{array}$$

$$x_t \leftarrow \nabla \varphi^* \left(-\eta \sum_{s < t} g_s \right)$$

$$\|g_t\|_* \leq G \text{ for all } t$$
$$\varphi : \mathcal{X} \to \mathbb{R} \text{ is 1-s.c. in } \|\cdot\|$$
$$\max_{x \in \mathcal{X}} \varphi(x) - \min_{x \in \mathcal{X}} \varphi(x) \leq \Theta$$

$$\max_{x^{\star} \in \mathcal{X}} \sum_{t \in [T]} \langle g_t, x_t - x^{\star} \rangle \lesssim G \sqrt{\Theta T}$$

$$x_t \leftarrow \nabla \varphi^* \left(-\eta \sum_{s < t} g_s \right)$$

= $\operatorname{argmin}_{x \in \mathcal{X}} \left\{ \eta \sum_{s < t} \langle g_s, x \rangle + \varphi(x) \right\}$

 $\|g_t\|_* \leq G \text{ for all } t$ $\varphi : \mathcal{X} \to \mathbb{R} \text{ is 1-s.c. in } \|\cdot\|$ $\max_{x \in \mathcal{X}} \varphi(x) - \min_{x \in \mathcal{X}} \varphi(x) \leq \Theta$

"follow the regularized leader"

$$\max_{x^{\star} \in \mathcal{X}} \sum_{t \in [T]} \langle g_t, x_t - x^{\star} \rangle \lesssim G\sqrt{\Theta T}$$

$$\|g_t\|_* \leq G \text{ for all } t$$
$$\varphi : \mathcal{X} \to \mathbb{R} \text{ is 1-s.c. in } \|\cdot\|$$
$$\max_{x \in \mathcal{X}} \varphi(x) - \min_{x \in \mathcal{X}} \varphi(x) \leq \Theta$$

 $\|g_t\|_* \leq G$ for all t

Gold standard (for εT regret over norm ball):

$$T = \operatorname{poly}\left(\frac{G\log(d)}{\epsilon}\right)$$



$$\mathcal{X} = \left\{ \mathbf{X} \in \mathrm{PSD}^{d \times d} \mid \mathrm{Tr}(\mathbf{X}) = 1 \right\}$$

Basic primitive: spectral bounds

$$\mathcal{X} = \left\{ \mathbf{X} \in \mathrm{PSD}^{d \times d} \mid \mathrm{Tr}(\mathbf{X}) = 1 \right\}$$

$$\varphi(\mathbf{X}) = \sum_{i \in [d]} \lambda_i(\mathbf{X}) \log \lambda_i(\mathbf{X})$$

Quantum (von Neumann) entropy

$$\mathcal{X} = \left\{ \mathbf{X} \in \mathrm{PSD}^{d \times d} \mid \mathrm{Tr}(\mathbf{X}) = 1 \right\}$$

$$\varphi(\mathbf{X}) = \sum_{i \in [d]} \lambda_i(\mathbf{X}) \log \lambda_i(\mathbf{X})$$

- Strongly convex?
- Bounded?
- Implementable?

$$\mathcal{X} = \left\{ \mathbf{X} \in \mathrm{PSD}^{d \times d} \mid \mathrm{Tr}(\mathbf{X}) = 1 \right\}$$

$$\varphi(\mathbf{X}) = \sum_{i \in [d]} \lambda_i(\mathbf{X}) \log \lambda_i(\mathbf{X})$$

Spoilers:

- Strongly convex in the trace norm
- Range: $\tilde{O}(1)$
- Matvec to $\nabla \varphi^*(\mathbf{Y})$ in near-linear time

$$\mathcal{X} = \left\{ \mathbf{X} \in \mathrm{PSD}^{d \times d} \mid \mathrm{Tr}(\mathbf{X}) = 1 \right\}$$

$$\varphi(\mathbf{X}) = \sum_{i \in [d]} \lambda_i(\mathbf{X}) \log \lambda_i(\mathbf{X})$$

Matrix multiplicative weights: mirror descent w.r.t. quantum entropy

Spoilers:

- Strongly convex in the trace norm
- Range: $\tilde{O}(1)$
- Matvec to ∇φ^{*}(Y) in near-linear time
$$\varphi^{*}(\mathbf{Y}) = \max_{\mathbf{X} \in \mathcal{X}} \langle \mathbf{Y}, \mathbf{X} \rangle - \sum_{i \in [d]} \lambda_{i}(\mathbf{X}) \log \lambda_{i}(\mathbf{X})$$
$$= \log \operatorname{Tr} \exp (\mathbf{Y})$$

$$\varphi^{*}(\mathbf{Y}) = \max_{\mathbf{X} \in \mathcal{X}} \langle \mathbf{Y}, \mathbf{X} \rangle - \sum_{i \in [d]} \lambda_{i}(\mathbf{X}) \log \lambda_{i}(\mathbf{X})$$
$$= \log \operatorname{Tr} \exp(\mathbf{Y})$$

Useful fact:
$$f_{ ext{mat}}(\mathbf{M}) = f_{ ext{vec}}(\lambda(\mathbf{M}))$$
 "spectral function" $\mathbf{M} = \mathbf{U} ext{diag}(\lambda(\mathbf{M}))\mathbf{U}^ op$

$$\varphi^{*}(\mathbf{Y}) = \max_{\mathbf{X} \in \mathcal{X}} \langle \mathbf{Y}, \mathbf{X} \rangle - \sum_{i \in [d]} \lambda_{i}(\mathbf{X}) \log \lambda_{i}(\mathbf{X})$$
$$= \log \operatorname{Tr} \exp(\mathbf{Y})$$

Useful fact:
$$f_{mat}(\mathbf{M}) = f_{vec}(\lambda(\mathbf{M}))$$
 convex, spectra
 $\nabla f_{mat}(\mathbf{X}) = \mathbf{U} \operatorname{diag}(\nabla f_{vec}(\lambda(\mathbf{M})))\mathbf{U}^{\top}$

$$\varphi^{*}(\mathbf{Y}) = \max_{\mathbf{X} \in \mathcal{X}} \langle \mathbf{Y}, \mathbf{X} \rangle - \sum_{i \in [d]} \lambda_{i}(\mathbf{X}) \log \lambda_{i}(\mathbf{X})$$
$$= \log \operatorname{Tr} \exp(\mathbf{Y})$$

Useful fact:
$$f_{\text{mat}}(\mathbf{M}) = f_{\text{vec}}(\lambda(\mathbf{M}))$$
 ...von Neumann!
 $\nabla f_{\text{mat}}(\mathbf{X}) = \mathbf{U} \text{diag}(\nabla f_{\text{vec}}(\lambda(\mathbf{M})))\mathbf{U}^{\top}$

$$\begin{split} \varphi^*(\mathbf{Y}) &= \max_{\mathbf{X} \in \mathcal{X}} \langle \mathbf{Y}, \mathbf{X} \rangle - \sum_{i \in [d]} \lambda_i(\mathbf{X}) \log \lambda_i(\mathbf{X}) \\ &= \log \operatorname{Tr} \exp \left(\mathbf{Y} \right) \\ \end{split}$$
Example: $f_{\operatorname{vec}}(\lambda) = \log \left(\sum_{i \in [d]} \exp(\lambda_i) \right) \quad \nabla f_{\operatorname{vec}}(\lambda) = \frac{\exp(\lambda)}{\|\exp(\lambda)\|_1}$

$$\begin{split} \varphi^*(\mathbf{Y}) &= \max_{\mathbf{X} \in \mathcal{X}} \langle \mathbf{Y}, \mathbf{X} \rangle - \sum_{i \in [d]} \lambda_i(\mathbf{X}) \log \lambda_i(\mathbf{X}) \\ &= \log \operatorname{Tr} \exp \left(\mathbf{Y} \right) \end{split}$$
Example:
$$\nabla \varphi^*(\mathbf{Y}) = \frac{\exp \left(\mathbf{Y} \right)}{\operatorname{Tr} \exp \left(\mathbf{Y} \right)}$$

Strong convexity

$$\varphi^*(\mathbf{Y}) = \log \operatorname{Tr} \exp(\mathbf{Y})$$

Useful fact: $v^{\top} \nabla^2 \varphi^*(y) v \le \|v\|_*^2$ $\iff v^{\top} \nabla^2 \varphi(x) v \ge \|v\|^2$

Strong convexity

$$\varphi^*(\mathbf{Y}) = \log \operatorname{Tr} \exp(\mathbf{Y})$$

Useful fact:

$$t: \quad v^{\top} \nabla^2 \varphi^*(y) v \le \|v\|_*^2$$

$$\iff v^{\top} \nabla^2 \varphi(x) v \ge \|v\|^2$$

"Smoothness-strong convexity duality"

Strong convexity

$$\varphi^*(\mathbf{Y}) = \log \operatorname{Tr} \exp(\mathbf{Y})$$

Useful fact:

$$t: \quad v^{\top} \nabla^2 \varphi^*(y) v \le \|v\|_*^2$$

$$\iff v^{\top} \nabla^2 \varphi(x) v \ge \|v\|^2$$

"Smoothness-strong convexity duality"

Why? Taylor expansion + $\varphi(x) = \frac{1}{2} \|x\|^2, \ \varphi^*(y) = \frac{1}{2} \|y\|_*^2$

Aside: "disentangling" lemma

$\operatorname{Tr}\left(\mathbf{M}^{\alpha}\mathbf{N}\mathbf{M}^{1-\alpha}\mathbf{N}\right) \leq \operatorname{Tr}(\mathbf{M}\mathbf{N}^{2})$

 $\mathbf{M}, \mathbf{N} \in \mathrm{PSD}^{d \times d}$ $\alpha \in [0, 1]$

Aside: "disentangling" lemma

$$\operatorname{Tr}\left(\mathbf{M}^{\alpha}\mathbf{N}\mathbf{M}^{1-\alpha}\mathbf{N}\right) \leq \operatorname{Tr}(\mathbf{M}\mathbf{N}^{2})$$

 $\mathbf{M}, \mathbf{N} \in \mathrm{PSD}^{d \times d}$ $\alpha \in [0, 1]$ Proof sketch:

$$\begin{pmatrix} \mathbf{N} & -\mathbf{N}^{\frac{1}{2}}\mathbf{M}^{\alpha}\mathbf{N}^{\frac{1}{2}} \\ -\mathbf{N}^{\frac{1}{2}}\mathbf{M}^{\alpha}\mathbf{N}^{\frac{1}{2}} & \mathbf{N}^{\frac{1}{2}}\mathbf{M}^{2\alpha}\mathbf{N}^{\frac{1}{2}} \end{pmatrix} \in \mathrm{PSD}^{2d \times 2d}$$

 $f(\alpha) = \operatorname{Tr}(\mathbf{M}^{\alpha}\mathbf{N}\mathbf{M}^{1-\alpha}\mathbf{N})$ is convex

$$\varphi^*(\mathbf{Y}) = \log \operatorname{Tr} \exp(\mathbf{Y})$$
$$\nabla^2 \varphi^*(\mathbf{Y})[\mathbf{M}, \mathbf{M}] \le \frac{1}{\operatorname{Tr} \exp(\mathbf{Y})} \langle \mathbf{M}, \nabla(\langle \mathbf{M}, \exp(\mathbf{Y}) \rangle) \rangle$$

$$\begin{split} \varphi^*(\mathbf{Y}) &= \log \operatorname{Tr} \exp(\mathbf{Y}) \\ \nabla^2 \varphi^*(\mathbf{Y})[\mathbf{M}, \mathbf{M}] \leq \frac{1}{\operatorname{Tr} \exp(\mathbf{Y})} \langle \mathbf{M}, \nabla(\langle \mathbf{M}, \exp(\mathbf{Y}) \rangle) \rangle \\ \langle \mathbf{M}, \nabla(\langle \mathbf{M}, \exp(\mathbf{Y}) \rangle) \rangle &= \sum_{k=1}^{\infty} \sum_{i=0}^{k-1} \frac{1}{k!} \langle \mathbf{M}, \mathbf{Y}^i \mathbf{M} \mathbf{Y}^{k-1-i} \rangle \end{split}$$

$$\begin{split} \varphi^*(\mathbf{Y}) &= \log \operatorname{Tr} \exp(\mathbf{Y}) \\ \nabla^2 \varphi^*(\mathbf{Y}) [\mathbf{M}, \mathbf{M}] \leq \frac{1}{\operatorname{Tr} \exp(\mathbf{Y})} \langle \mathbf{M}, \nabla(\langle \mathbf{M}, \exp(\mathbf{Y}) \rangle) \rangle \\ \langle \mathbf{M}, \nabla(\langle \mathbf{M}, \exp(\mathbf{Y}) \rangle) \rangle &= \sum_{k=1}^{\infty} \sum_{i=0}^{k-1} \frac{1}{k!} \langle \mathbf{M}, \mathbf{Y}^i \mathbf{M} \mathbf{Y}^{k-1-i} \rangle \\ &\leq \sum_{k=0}^{\infty} \frac{1}{k!} \langle \mathbf{M}^2, \mathbf{Y}^k \rangle \quad \text{(disentangling)} \end{split}$$

$$\begin{split} \varphi^*(\mathbf{Y}) &= \log \operatorname{Tr} \exp(\mathbf{Y}) \\ \nabla^2 \varphi^*(\mathbf{Y})[\mathbf{M}, \mathbf{M}] \leq \frac{1}{\operatorname{Tr} \exp(\mathbf{Y})} \langle \mathbf{M}, \nabla(\langle \mathbf{M}, \exp(\mathbf{Y}) \rangle) \rangle \\ \langle \mathbf{M}, \nabla(\langle \mathbf{M}, \exp(\mathbf{Y}) \rangle) \rangle &= \sum_{k=1}^{\infty} \sum_{i=0}^{k-1} \frac{1}{k!} \langle \mathbf{M}, \mathbf{Y}^i \mathbf{M} \mathbf{Y}^{k-1-i} \rangle \\ &\leq \sum_{k=0}^{\infty} \frac{1}{k!} \langle \mathbf{M}^2, \mathbf{Y}^k \rangle = \langle \mathbf{M}^2, \exp(\mathbf{Y}) \rangle \end{split}$$

$$\varphi^*(\mathbf{Y}) = \log \operatorname{Tr} \exp(\mathbf{Y})$$
$$\nabla^2 \varphi^*(\mathbf{Y})[\mathbf{M}, \mathbf{M}] \le \left\langle \mathbf{M}^2, \frac{\exp(\mathbf{Y})}{\operatorname{Tr} \exp(\mathbf{Y})} \right\rangle$$

$$\begin{split} \varphi^*(\mathbf{Y}) &= \log \operatorname{Tr} \exp(\mathbf{Y}) \\ \nabla^2 \varphi^*(\mathbf{Y})[\mathbf{M}, \mathbf{M}] \leq \left\langle \mathbf{M}^2, \frac{\exp(\mathbf{Y})}{\operatorname{Tr} \exp(\mathbf{Y})} \right\rangle \\ &\leq \lambda_{\max}(\mathbf{M}^2) = \left\| \mathbf{M} \right\|_{\infty}^2 \end{split}$$

$$\begin{split} \varphi^*(\mathbf{Y}) &= \log \operatorname{Tr} \exp(\mathbf{Y}) \\ \nabla^2 \varphi^*(\mathbf{Y})[\mathbf{M}, \mathbf{M}] \leq \left\langle \mathbf{M}^2, \frac{\exp(\mathbf{Y})}{\operatorname{Tr} \exp(\mathbf{Y})} \right\rangle \\ &\leq \lambda_{\max}(\mathbf{M}^2) = \|\mathbf{M}\|_{\infty}^2 \end{split}$$

....so, entropy is strongly convex in Schatten-1!

$$\begin{split} \varphi^*(\mathbf{Y}) &= \log \operatorname{Tr} \exp(\mathbf{Y}) \\ \nabla^2 \varphi^*(\mathbf{Y})[\mathbf{M}, \mathbf{M}] \leq \left\langle \mathbf{M}^2, \frac{\exp(\mathbf{Y})}{\operatorname{Tr} \exp(\mathbf{Y})} \right\rangle \\ &\leq \lambda_{\max}(\mathbf{M}^2) = \left\| \mathbf{M} \right\|_{\infty}^2 \end{split}$$

"local norms" smoothness bound

Implementation

 $\mathbf{Y} = -\eta \sum \mathbf{G}_s$ s < t $\nabla \varphi^*(\mathbf{Y}) = \frac{\exp \mathbf{Y}}{\operatorname{Tr} \exp(\mathbf{Y})}$

Implementation

$$\mathbf{Y} = -\eta \sum_{s < t} \mathbf{G}_s \qquad \|\mathbf{Y}\|_{\text{op}} = \text{poly}\left(\frac{G \log(d)}{\epsilon}\right)$$

$$\nabla \varphi^*(\mathbf{Y}) = \frac{\exp \mathbf{Y}}{\operatorname{Tr} \exp(\mathbf{Y})}$$

Our action: need to "access" efficiently Implementation

 $a \approx v^{\top} \exp(\mathbf{Y}) v$

Warmup: single quadratic form

Implementation

 $a \approx v^{\top} \exp(\mathbf{Y}) v$

Key idea: polynomial approximation

degree- Δp

 $\implies \mathcal{T}_{\mathrm{mv}}(p(\mathbf{Y})) = O(\mathcal{T}_{\mathrm{mv}}(\mathbf{Y}) \cdot \Delta)$

Implementation

 $a \approx v^{\top} \exp(\mathbf{Y}) v$

$p(x) \approx \exp(x)$ for $x \in [\lambda_{\min}(\mathbf{Y}), \lambda_{\max}(\mathbf{Y})]$

Implementation

 $a \approx v^{\top} \exp(\mathbf{Y})v$

 $p(x) \approx \exp(x)$ for $x \in [\lambda_{\min}(\mathbf{Y}), \lambda_{\max}(\mathbf{Y})]$

 $\lambda_{\max}(\mathbf{M}) - \lambda_{\min}(\mathbf{M}) \le R$

Degree $\approx R$ Taylor approximation is high-accuracy

Implementation

 $a = v^{\top} p(\mathbf{Y}) v$

Computable in "nearly-linear time":

$$\mathcal{T}_{\mathrm{mv}}(\mathbf{Y}) \cdot \mathrm{poly}\left(\frac{G\log(d)}{\epsilon}\right)$$

Implementation

$$\left\{a_i \approx v_i^\top \exp\left(\mathbf{Y}\right) v_i\right\}_{i \in [n]}$$

...what good is one quadratic form?

Implementation

$$\left\{a_i \approx v_i^\top \exp\left(\mathbf{Y}\right) v_i\right\}_{i \in [n]}$$

$$v^{\top} \exp(\mathbf{Y}) v \approx v^{\top} \mathbf{Q} \exp(\mathbf{Y}) \mathbf{Q}^{\top} v$$

(**Q** is any JL matrix)

Key idea: reuse multiplies via sketching (warning: independence)

Implementation

$$\left\{a_i \approx v_i^\top \exp\left(\mathbf{Y}\right) v_i\right\}_{i \in [n]}$$

$$\begin{pmatrix} \tilde{q}_1^\top \\ \vdots \\ \tilde{q}_k^\top \end{pmatrix} = p \begin{pmatrix} \frac{1}{2} \mathbf{Y} \end{pmatrix} \begin{pmatrix} q_1^\top \\ \vdots \\ q_k^\top \end{pmatrix}$$

Implementation

$$\left\{a_i \approx v_i^\top \exp\left(\mathbf{Y}\right) v_i\right\}_{i \in [n]}$$



Implementation

$$\left\{a_i \approx v_i^\top \exp\left(\mathbf{Y}\right) v_i\right\}_{i \in [n]}$$

$$\dots \text{at ``test time''} \dots$$
$$a_i = \left\| \begin{pmatrix} \tilde{q}_1^\top \\ \vdots \\ \tilde{q}_k^\top \end{pmatrix} v_i \right\|_2^2$$

runtime:

$$d \cdot \operatorname{poly}\left(\frac{G\log(d)}{\epsilon}\right)$$

$$k = \operatorname{poly}\left(\frac{G\log(d)}{\epsilon}\right)$$

MMW summary

$$\max_{\substack{\mathbf{X}^{\star} \in \mathrm{PSD}^{d \times d} \\ \mathrm{Tr}(\mathbf{X}^{\star})=1}} \sum_{t \in [T]} \langle \mathbf{G}_{t}, \mathbf{X}^{\star} - \mathbf{X}_{t} \rangle \lesssim G\sqrt{T}$$

$$\|\mathbf{G}_t\|_{\infty} \leq G \text{ for all } t \in [T]$$

MMW summary

$$\max_{\substack{\mathbf{X}^{\star} \in \mathrm{PSD}^{d \times d} \\ \mathrm{Tr}(\mathbf{X}^{\star}) = 1}} \frac{1}{T} \sum_{t \in [T]} \langle \mathbf{G}_{t}, \mathbf{X}^{\star} - \mathbf{X}_{t} \rangle \leq \epsilon$$

$$T = \operatorname{poly}\left(\frac{G\log(d)}{\epsilon}\right)$$

iteration count

$$\max_{\substack{\mathbf{X}^{\star} \in \mathrm{PSD}^{d \times d} \\ \mathrm{Tr}(\mathbf{X}^{\star}) = 1}} \frac{1}{T} \sum_{t \in [T]} \langle \mathbf{G}_{t}, \mathbf{X}^{\star} - \mathbf{X}_{t} \rangle \leq \epsilon$$

$$T = \operatorname{poly}\left(\frac{G \log(d)}{\epsilon}\right) \quad \left(\sum_{t \in [T]} \mathcal{T}_{\operatorname{mv}}(\mathbf{G}_t)\right) \cdot \operatorname{poly}\left(\frac{G \log(d)}{\epsilon}\right)$$

iteration count cost for "implementing" each iteration

Improvement: runtime

$$\max_{\substack{\mathbf{X}^{\star} \in \mathrm{PSD}^{d \times d} \\ \mathrm{Tr}(\mathbf{X}^{\star})=1}} \frac{1}{T} \sum_{t \in [T]} \langle \mathbf{G}_{t}, \mathbf{X}^{\star} - \mathbf{X}_{t} \rangle \leq \epsilon$$

[CDST19], see also [BBN13] for different SOTA tradeoff

 $\left(\frac{G\log(d)}{\epsilon}\right)^2$

iteration count

 $\left(\frac{G\log(d)}{\epsilon}\right)^{0.5}$

cost for "implementing" each iteration

Improvement: local norms

"prediction error" per iteration, controlled by s.c.

$$\max_{\substack{\mathbf{X}^{\star} \in \mathrm{PSD}^{d \times d} \\ \mathrm{Tr}(\mathbf{X}^{\star})=1}} \sum_{t \in [T]} \langle \mathbf{G}_{t}, \mathbf{X}^{\star} - \mathbf{X}_{t} \rangle \lesssim \frac{1}{\eta} + \eta G^{2} T$$

size of regularizer
Improvement: local norms

$$\max_{\substack{\mathbf{X}^{\star} \in \mathrm{PSD}^{d \times d} \\ \mathrm{Tr}(\mathbf{X}^{\star})=1}} \sum_{t \in [T]} \langle \mathbf{G}_{t}, \mathbf{X}^{\star} - \mathbf{X}_{t} \rangle \lesssim \frac{1}{\eta} + \eta G^{2}T$$

$$\max_{\substack{\mathbf{X}^{\star} \in \mathrm{PSD}^{d \times d} \\ \mathrm{Tr}(\mathbf{X}^{\star})=1}} \sum_{t \in [T]} \langle \mathbf{G}_{t}, \mathbf{X}^{\star} - \mathbf{X}_{t} \rangle \lesssim \frac{1}{\eta} + \eta G \sum_{t \in [T]} \langle \mathbf{G}_{t}, \mathbf{X}_{t} \rangle$$

can drastically improve if \mathbf{G}_t reacts to \mathbf{X}_t

Extension: Schatten-norm setups

$$\varphi(\mathbf{X}) = \frac{1}{2(q-1)} \|\mathbf{X}\|_q^2$$

globally I-s.c. in Schatten-q norm

$$\varphi(\mathbf{X}) = \frac{1}{2q(q-1)} \|\mathbf{X}\|_q^q$$

I-s.c. in Schatten-q norm on unit ball

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I-s.c. in Schatten-*q* norm on unit ball

Who cares? ...better captures multiplicative (vs. additive) ...offers different tradeoffs (e.g. lower moment bounds)

Extension: Positive SDP

 $\min_{\mathbf{X}\in\mathcal{X}}\max_{y\in\mathcal{Y}}\left\langle \mathbf{X},\sum_{i\in[n]}y_i\mathbf{A}_i\right\rangle$

Canonical application: feasibility SDP via saddle points

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Canonical application: feasibility SDP via saddle points

If all A_i are PSD or NSD, can get *multiplicative* error guarantees with no dependence on "width"

$$G := \max_{i \in [n]} \lambda_{\max} \left(\mathbf{A}_i \right)$$

Extension: Positive SDP

 $\min_{\mathbf{X}\in\mathcal{X}}\max_{y\in\mathcal{Y}}\left\langle \mathbf{X},\sum_{i\in[n]}y_i\mathbf{A}_i\right\rangle$

Canonical application: feasibility SDP via saddle points

If all A_i are PSD or NSD, can get *multiplicative* error guarantees with no dependence on "width"

- Works at every scale
- Often the case in robust statistics! (Sample covariances)

Talk outline

- A gentle introduction to MMW
 - Regret minimization
 - Matrix analysis
 - Implementation
 - Relatives of MMW

• Robust statistics primitives via MMW

- Mean estimation
- A tour of applications





$$\begin{split} \{X_i^{\star}\}_{i\in[n]} \sim_{\text{i.i.d.}} \mathcal{D} \\ \{X_i = X_i^{\star}\}_{i\in G} \\ \{X_i\}_{i\in B}, \ |B| \approx \epsilon n \end{split}$$

Goal: estimate $\mu^{\star} = \mu(\mathcal{D})$

Setting

 $\frac{1}{|G|} \sum_{i \in G} (X_i - \mu^*) (X_i - \mu^*)^\top \preceq \mathbf{I}_d$ Meta-algo

$$\{X_i^{\star}\}_{i\in[n]} \sim_{\text{i.i.d.}} \mathcal{D}$$
$$\{X_i = X_i^{\star}\}_{i\in G}$$
$$\{X_i\}_{i\in B}, |B| \approx \epsilon n$$

Goal: estimate $\ \mu^{\star} = \mu(\mathcal{D})$

Setting

 $\frac{1}{|G|} \sum_{i \in C} (X_i - \mu^\star) (X_i - \mu^\star)^\top \preceq \mathbf{I}_d$ Return empirical mean μ_w if: $\sum w_i (X_i - \mu_w) (X_i - \mu_w)^\top \preceq O(1) \mathbf{I}_d$ $i \in [n]$ $\sum w_i \le 1, \ \sum w_i \ge 1 - O(\epsilon)$ $i \in [n]$ $i \in G$ Meta-algo

$$\begin{split} \{X_i^{\star}\}_{i\in[n]} \sim_{\text{i.i.d.}} \mathcal{D} \\ \{X_i = X_i^{\star}\}_{i\in G} \\ \{X_i\}_{i\in B}, \ |B| \approx \epsilon n \end{split}$$

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Setting

 $\frac{1}{|G|} \sum_{i \in G} (X_i - \mu^\star) (X_i - \mu^\star)^\top \preceq \mathbf{I}_d$ Return empirical mean μ_w if: $\sum w_i (X_i - \mu_w) (X_i - \mu_w)^\top \preceq O(1) \mathbf{I}_d$ $i \in [n]$ $\sum w_i \le 1, \ \sum w_i \ge 1 - O(\epsilon)$ $i \in [n]$ $i \in G$ Invariant: Meta-algo "saturation"

$$\begin{split} \{X_i^{\star}\}_{i\in[n]} \sim_{\text{i.i.d.}} \mathcal{D} \\ \{X_i = X_i^{\star}\}_{i\in G} \\ \{X_i\}_{i\in B}, \ |B| \approx \epsilon n \\ \text{Goal: estimate} \ \mu^{\star} = \mu(\mathcal{D}) \\ \text{Setting} \end{split}$$

$$\frac{1}{|G|} \sum_{i \in G} (X_i - \mu^*) (X_i - \mu^*)^\top \preceq \mathbf{I}_d$$

Else:
$$\exists \mathbf{X} \in \mathrm{PSD}^{d \times d} : \mathrm{Tr} \mathbf{X} = 1$$

$$\mathbb{E}_{i \sim w} \left\langle (X_i - \mu_w) (X_i - \mu_w)^\top, \mathbf{X} \right\rangle \gg 1$$

Meta-algo

$$\{X_i^{\star}\}_{i\in[n]} \sim_{\text{i.i.d.}} \mathcal{D}$$
$$\{X_i = X_i^{\star}\}_{i\in G}$$
$$\{X_i\}_{i\in B}, |B| \approx \epsilon n$$

Goal: estimate $\ \mu^{\star} = \mu(\mathcal{D})$

Setting

$$\frac{1}{|G|} \sum_{i \in G} (X_i - \mu^*) (X_i - \mu^*)^\top \preceq \mathbf{I}_d$$
Else:

$$\exists \mathbf{X} \in \mathrm{PSD}^{d \times d} : \mathrm{Tr} \mathbf{X} = 1$$

$$\mathbb{E}_{i \sim w} \left\langle (X_i - \mu_w) (X_i - \mu_w)^\top, \mathbf{X} \right\rangle \gg 1$$
Many fast ways of preserving saturation
Meta-algo





$$\max_{\substack{\mathbf{X}^{\star} \in \mathrm{PSD}^{d \times d} \\ \mathrm{Tr}(\mathbf{X}^{\star})=1}} \sum_{t \in [T]} \langle \mathbf{G}_{t}, \mathbf{X}^{\star} - \mathbf{X}_{t} \rangle \lesssim \frac{1}{\eta} + \eta G \sum_{t \in [T]} \langle \mathbf{G}_{t}, \mathbf{X}_{t} \rangle$$

$$\max_{\substack{\mathbf{X}^{\star} \in \mathrm{PSD}^{d \times d} \\ \mathrm{Tr}(\mathbf{X}^{\star})=1}} \sum_{t \in [T]} \langle \mathbf{G}_{t}, \mathbf{X}^{\star} - \mathbf{X}_{t} \rangle \lesssim \frac{1}{\eta} + \eta G \sum_{t \in [T]} \langle \mathbf{G}_{t}, \mathbf{X}_{t} \rangle$$

$$\left\| \sum_{t \in [T]} \mathbf{G}_t \right\|_{\text{op}} \lesssim G + 2 \sum_{t \in [T]} \langle \mathbf{G}_t, \mathbf{X}_t \rangle$$

$$\mathbf{G}_t = \sum_{i \in [n]} [w_t]_i \left(X_i - \mu_{w_t} \right) \left(X_i - \mu_{w_t} \right)^\top \qquad \mathbf{G}_0 \preceq G_0 \mathbf{I}_d$$
$$G_0 \gg 1$$

$$\left\| \sum_{t \in [T]} \mathbf{G}_t \right\|_{\text{op}} \lesssim G_0 + 2 \sum_{t \in [T]} \langle \mathbf{G}_t, \mathbf{X}_t \rangle$$

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$$G_0 \gg 1$$

$$\left\| \sum_{t \in [T]} \mathbf{G}_t \right\|_{\text{op}} \lesssim G_0 + 2 \sum_{t \in [T]} \langle \mathbf{G}_t, \mathbf{X}_t \rangle$$

 $\lesssim G_0 + O(T)$ (filter in each iteration)

$$\mathbf{G}_t = \sum_{i \in [n]} [w_t]_i \left(X_i - \mu_{w_t} \right) \left(X_i - \mu_{w_t} \right)^\top \qquad \mathbf{G}_0 \preceq G_0 \mathbf{I}_d$$
$$G_0 \gg 1$$

$$T \|\mathbf{G}_T\|_{\mathrm{op}} \lesssim \left\| \sum_{t \in [T]} \mathbf{G}_t \right\|_{\mathrm{op}} \lesssim G_0 + 2 \sum_{t \in [T]} \langle \mathbf{G}_t, \mathbf{X}_t \rangle$$

 $\begin{array}{ll} \mbox{monotone} & \qquad \qquad \lesssim G_0 + O(T) \\ \mbox{feedbacks} & \qquad \qquad \end{array}$

$$\mathbf{G}_t = \sum_{i \in [n]} [w_t]_i \left(X_i - \mu_{w_t} \right) \left(X_i - \mu_{w_t} \right)^\top \quad \mathbf{G}_0 \preceq \mathbf{G}_0 \mathbf{I}_d$$

$$T \|\mathbf{G}_{T}\|_{\mathrm{op}} \lesssim \left\| \sum_{t \in [T]} \mathbf{G}_{t} \right\|_{\mathrm{op}} \lesssim G_{0} + 2 \sum_{t \in [T]} \langle \mathbf{G}_{t}, \mathbf{X}_{t} \rangle \qquad T \lesssim 1$$

monotone
feedbacks
$$\lesssim G_{0} + O(T) \qquad \|\mathbf{G}_{T}\|_{\mathrm{op}} \leq \frac{G_{0}}{2}$$
Punchline

$$\mathbf{G}_t = \sum_{i \in [n]} [w_t]_i \left(X_i - \mu_{w_t} \right) \left(X_i - \mu_{w_t} \right)^\top \quad \mathbf{G}_0 \preceq G_0 \mathbf{I}_d$$

Interpretation:

MMW as a multidirectional filter



$$\mathbf{G}_t = \sum_{i \in [n]} [w_t]_i \left(X_i - \mu_{w_t} \right) \left(X_i - \mu_{w_t} \right)^\top \quad \mathbf{G}_0 \preceq \mathbf{G}_0 \mathbf{I}_d$$

Interpretation:

MMW as a multidirectional filter

...[DHL '19] Robust mean estimation in time $\tilde{O}(nd)$



$\operatorname{Tr} (\mathbf{Y}^p) \propto \max_{\substack{\mathbf{X} \in \operatorname{Sym}^{d \times d} \\ \|\mathbf{X}\|_q \leq 1}} \langle \mathbf{X}, \mathbf{Y} \rangle$

as a potential

$\operatorname{Tr}(\mathbf{Y}^{p}) \propto \max_{\substack{\mathbf{X} \in \operatorname{Sym}^{d \times d} \\ \|\mathbf{X}\|_{q} \leq 1}} \langle \mathbf{X}, \mathbf{Y} \rangle$

as a potential

Upshot:

- Single-iteration progress (MMW non-monotone)
 - Multifilter [DKKLT '22], list-decoding

$\operatorname{Tr}(\mathbf{Y}^{p}) \propto \max_{\substack{\mathbf{X} \in \operatorname{Sym}^{d \times d} \\ \|\mathbf{X}\|_{q} \leq 1}} \langle \mathbf{X}, \mathbf{Y} \rangle$

as a potential

Upshot:

- Single-iteration progress (MMW non-monotone)
 - Multifilter [DKKLT '22], list-decoding
- More natural interpretation?
 - Power method [DKKP '23], PCA

$\operatorname{Tr}(\mathbf{Y}^{p}) \propto \max_{\substack{\mathbf{X} \in \operatorname{Sym}^{d \times d} \\ \|\mathbf{X}\|_{q} \leq 1}} \langle \mathbf{X}, \mathbf{Y} \rangle$

as a potential

Downside(?)

- Less obvious connection to regret minimization
- (Does not apply to Daniel Kane)

$\operatorname{Tr}(\mathbf{Y}^{p}) \propto \max_{\substack{\mathbf{X} \in \operatorname{Sym}^{d \times d} \\ \|\mathbf{X}\|_{q} \leq 1}} \langle \mathbf{X}, \mathbf{Y} \rangle$

as a potential

Downside(?)

- Less obvious connection to regret minimization
- Suggest: mirror descent as a catch-all
- Smarter filters for specific problem

$$\min_{\substack{w \in \mathbb{R}^n_{\geq 0} \\ \|w\|_1 \leq 1}} \left\| \sum_{i \in [n]} w_i \left(X_i - \mu_w \right) \left(X_i - \mu_w \right)^\top \right\|_{\text{op}}$$

Packing SDP

$$\min_{\substack{w \in \mathbb{R}^n_{\geq 0} \\ \|w\|_1 \leq 1}} \left\| \sum_{i \in [n]} w_i \left(X_i - \mu_w \right) \left(X_i - \mu_w \right)^\top \right\|_{\text{op}}$$

Use case: local reweightings (e.g. gradient descent)

Iterative methods: O(I) approx. is OK [PSBR '18, CDG '19, ...]

$$\min_{\substack{w \in \mathbb{R}^n_{\geq 0} \\ \|w\|_1 \leq 1}} \left\| \sum_{i \in [n]} w_i \left(X_i - \mu_w \right) \left(X_i - \mu_w \right)^\top \right\|_{\text{op}}$$



value = step size dual = descent direction

Very general strategy for stochastic optimization problems...

$$\mathcal{X} = \left\{ \mathbf{X} \in \mathrm{PSD}^{d \times d} \mid \|\mathbf{X}\|_{\mathrm{op}} \le 1, \mathrm{TrX} \le k \right\}$$

"Fantope" = cvx hull of projection matrices

$$\mathcal{X} = \left\{ \mathbf{X} \in \mathrm{PSD}^{d \times d} \mid \|\mathbf{X}\|_{\mathrm{op}} \le 1, \mathrm{TrX} \le k \right\}$$

"Fantope" = cvx hull of projection matrices Multi-direction filters: list-decoding [DKKLT '21], optimal Huber contamination [DKPP '23]

$$\mathbf{G} = \frac{1}{\kappa} \sum_{i \in [n]} w_i \mathbf{A}_i$$

$$\sum_{i\in[n]} w_i \mathbf{A}_i \preceq \mathbf{I}_d$$

e.g. solution to a packing SDP

$$\mathbf{G} = \frac{1}{\kappa} \sum_{i \in [n]} w_i \mathbf{A}_i$$

$$\frac{O(1)}{\kappa} \mathbf{I}_d \preceq \sum_{i \in [n]} \bar{w}_i \mathbf{A}_i \preceq \mathbf{I}_d$$

Regret minimization: two-sided constraints

 $\sum w_i \mathbf{A}_i \preceq \mathbf{I}_d$ $i \in [n]$

$$\mathbf{G} = \frac{1}{\kappa} \sum_{i \in [n]} w_i \mathbf{A}_i$$

 $\sum w_i \mathbf{A}_i \preceq \mathbf{I}_d$

 $i \in [n]$

$$\frac{O(1)}{\kappa} \mathbf{I}_d \preceq \sum_{i \in [n]} \bar{w}_i \mathbf{A}_i \preceq \mathbf{I}_d$$

Regret minimization: two-sided constraints

e.g. planted well-conditioning, semi-random linear models [JLMSST '23]
Thank you!

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